SPREADING IN CLAW-FREE CUBIC GRAPHS

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Communicated by Dalibor Fronček

Abstract. Let $p \in \mathbb{N}$ and $q \in \mathbb{N} \cup \{\infty\}$. We study a dynamic coloring of the vertices of a graph G that starts with an initial subset S of blue vertices, with all remaining vertices colored white. If a white vertex v has at least p blue neighbors and at least one of these blue neighbors of v has at most q white neighbors, then by the spreading color change rule the vertex v is recolored blue. The initial set S of blue vertices is a (p,q)-spreading set for G if by repeatedly applying the spreading color change rule all the vertices of G are eventually colored blue. The (p,q)-spreading set is a generalization of the well-studied concepts of k-forcing and r-percolating sets in graphs. For $q \geq 2$, a (1,q)-spreading set is exactly a q-forcing set, and the (1,1)-spreading set is a 1-forcing set (also called a zero forcing set), while for $q=\infty$, a (p, ∞) -spreading set is exactly a p-percolating set. The (p, q)-spreading number, $\sigma_{(p,q)}(G)$, of G is the minimum cardinality of a (p,q)-spreading set. In this paper, we study (p,q)-spreading in claw-free cubic graphs. While the zero-forcing number of claw-free cubic graphs was studied earlier, for each pair of values p and q that are not both 1 we either determine the (p,q)-spreading number of a claw-free cubic graph G or show that $\sigma_{(p,q)}(G)$ attains one of two possible values.

Keywords: bootstrap percolation, zero forcing set, k-forcing set, spreading.

Mathematics Subject Classification: 05C35, 05C75.

1. INTRODUCTION

In this paper we continue the study of dynamic graph colorings and explore the concept of (p,q)-spreading in claw-free cubic graph. The concept of spreading in graphs is a generalization of the well-studied concepts of k-forcing and r-percolating sets in graphs.

Consider a dynamic coloring of the vertices of a graph G that starts with an initial subset S of blue vertices, with all remaining vertices colored white. For $k \in \mathbb{N}$, the color change rule in k-forcing is defined as follows: if a blue vertex u has at most k

white neighbors, then all white neighbors of u are recolored blue. In particular, when k=1, this is the color change rule for zero forcing in graphs. The initial set S of blue vertices is a k-forcing set for G if by repeatedly applying the color change rule in k-forcing all the vertices of G are eventually colored blue. The k-forcing number of G, denoted $F_k(G)$, is the minimum cardinality among all k-forcing set of G. When k=1, the k-forcing number of G is called the zero forcing number and is denoted by Z(G), and so $Z(G) = F_1(G)$.

The concept of bootstrap percolation has a similar definition to that of zero forcing, yet a different motivation. It was introduced as a simplified model of a magnetic system in 1979 [10], and was later studied on random graphs and also on deterministic graphs. In particular, already some early papers considered bootstrap percolation in grids [7, 8].

As before, we consider a dynamic coloring of the vertices of a graph G that starts with an initial subset S of blue vertices, with all remaining vertices colored white. We refer to blue vertices as "infected" and white vertices as "uninfected". For $r \in \mathbb{N}$, the color change rule in r-percolation is defined as follows: if a white (uninfected) vertex u has at least r blue (infected) neighbors, then the vertex u is recolored blue. The initial set S of blue vertices is an r-neighbor bootstrap percolating set, or simply an r-percolating set, of G if by repeatedly applying the color change rule in r-percolation all the vertices of G are eventually colored blue. The r-neighbor bootstrap percolation number, or simply the r-percolation number, of G, denoted m(G,r), is the minimum cardinality among all r-neighbor bootstrap percolating sets of G.

Motivated by the similarity of the definitions of the above two concepts, a common generalization of k-forcing and r-bootstrap percolation was introduced in [9]. Let $p \in \mathbb{N}$ and $q \in \mathbb{N} \cup \{\infty\}$. As before, we consider a dynamic coloring of the vertices of a graph G that starts with an initial subset S of blue vertices, with all remaining vertices colored white. The color change rule in (p,q)-spreading is defined as follows: if a white vertex w has at least p blue neighbors, and one of the blue neighbors of w has at most q white neighbors, then the vertex w is recolored blue. The initial set S of blue vertices is a (p,q)-spreading set of G, if by repeatedly applying the (p,q)-spreading color change rule all the vertices of G are eventually colored blue. The (p,q)-spreading number, denoted $\sigma_{(p,q)}(G)$, of a graph G is the minimum cardinality among all (p,q)-spreading sets.

The (p,q)-spreading set is a generalization of the well-studied concepts of k-forcing and r-percolating sets in graphs. For $q \geq 2$, a (1,q)-spreading set is exactly a q-forcing set [3], and the (1,1)-spreading set is a 1-forcing set (also called a zero forcing set [2,21]), while if $q \geq \Delta(G)$ (including $q=\infty$), then a (p,q)-spreading set is exactly a p-percolating set. In the foundational paper [9], the complexity of the decision version of the (p,q)-spreading number was proved to be NP-complete for all p and q, while efficient algorithms for determining these numbers were found in trees. In addition, for almost all values of p and q, the (p,q)-spreading numbers of Cartesian grids $P_n \square P_m$ were established in [9].

The following trivial observation will be (at least implicitly) used several times in the paper.

Observation 1.1. Let G be a graph, $p \ge 2$ and $q \in \mathbb{N} \cup \{\infty\}$. If P is a (p,q)-spreading set in G, then P is also a (p,q+1)-spreading set in G as well as a (p-1,q)-spreading set in G. In particular,

$$\sigma_{(p,q)}(G) \geq \sigma_{(p,q+1)}(G) \ \ and \ \ \sigma_{(p,q)}(G) \geq \sigma_{(p-1,q)}(G).$$

A graph is *claw-free* if it does not contain $K_{1,3}$ as an induced subgraph. A *cubic graph* (also called a 3-*regular graph*) is a graph in which every vertex has degree 3. In this paper we study spreading in claw-free cubic graphs.

1.1. GRAPH THEORY NOTATION AND TERMINOLOGY

For notation and graph theory terminology, we in general follow [17]. Specifically, let G be a graph with vertex set V(G) and edge set E(G), and of order n(G) = |V(G)| and size m(G) = |E(G)|. A neighbor of a vertex v in G is a vertex u that is adjacent to v, that is, $uv \in E(G)$. The open neighborhood $N_G(v)$ of a vertex v in G is the set of neighbors of v, while the closed neighborhood of v is the set $N_G[v] = \{v\} \cup N(v)$. We denote the degree of v in G by $\deg_G(v) = |N_G(v)|$, and $\Delta(G) = \max\{\deg_G(v) : v \in V(G)\}$. For a set $S \subseteq V(G)$, its open neighborhood is the set $N_G(S) = \bigcup_{v \in S} N_G(v)$, and its closed neighborhood is the set $N_G[S] = N_G(S) \cup S$.

For a set $S \subseteq V(G)$, the subgraph induced by S is denoted by G[S]. Further, the subgraph of G obtained from G by deleting all vertices in S and all edges incident with vertices in S is denoted by G - S; that is, $G - S = G[V(G) \setminus S]$. If $S = \{v\}$, then we also denote G - S simply by G - v. If F is a graph, then an F-component of G is a component isomorphic to F. We denote by $\alpha(G)$ the independence number of G, and so $\alpha(G)$ is the maximum cardinality among all independent sets of G. A vertex cover of G is a set S of vertices such that every edge of G is incident with at least one vertex in S. The vertex covering number $\beta(G)$ (also denoted $\tau(G)$ in the literature), equals the minimum cardinality among all vertex covers of G. By the well-known Gallai theorem, $\alpha(G) + \beta(G) = n(G)$ holds in any graph G.

We denote the path, cycle, and complete graph on n vertices by P_n , C_n , and K_n , respectively, and we denote the complete bipartite graph with partite sets of cardinality n and m by $K_{n,m}$. A triangle in G is a subgraph isomorphic to K_3 , whereas a diamond in G is an induced subgraph of G isomorphic to K_4 with one edge missing, denoted by $K_4 - e$. A graph is diamond-free if it does not contain $K_4 - e$ as an induced subgraph. For $k \geq 1$ an integer, we use the standard notation $[k] = \{1, \ldots, k\}$.

1.2. MAIN RESULTS AND ORGANIZATION OF THE PAPER

Davila and Henning studied zero forcing in claw-free cubic graphs [11]. They established two upper bounds, which are summarized in the following result (for the definition of diamond-necklace N_k , see Section 2).

Theorem 1.2 ([11]). If $G \neq K_4$ is a connected, claw-free, cubic graph of order n, then the following properties hold:

- (a) $Z(G) \le \alpha(G) + 1$,
- (b) $Z(G) \le \frac{n}{3} + 1$, unless $G = N_2$ in which case $Z(G) = \frac{1}{3}(n+4)$.

It is proved in [11] that both upper bounds in Theorem 1.2 are sharp. He et al. [18] recently characterized the connected, claw-free cubic graphs G for which $Z(G) = \alpha(G) + 1$. Notably, there are only three sporadic graphs that attain this value, namely N_2 , N_3 and the Hamming graph $K_3 \square K_2$. Since these graphs have 8,12 and 6 vertices, respectively, the following result immediately follows.

Theorem 1.3 ([18]). If $G \neq K_4$ is a connected, claw-free, cubic graph of order at least 14, then $Z(G) \leq \alpha(G)$.

Note that $Z(G) = \sigma_{(1,1)}(G)$, and the mentioned result (Theorem 1.2(a)) is placed in the (1,1)-entry of Table 1.

From properties of connected, claw-free, cubic graphs G which we discuss in Section 2, if $G \neq K_4$ then there is a unique partition of the vertex set V(G) into subsets each of which induces either a triangle or a diamond. The number of these subsets, called units, in G is denoted by u(G). Table 1 summarizes the values of (p,q)-spreading numbers in claw-free cubic graphs $G \neq K_4$ for all possible p and q in terms of the parameters $\alpha(G)$, $\beta(G)$, u(G) and n(G). Besides the (1,1)-entry giving the mentioned upper bound on $\sigma_{(1,1)}(G)$ all other values in Table 1 are obtained in this paper. It is easy to see that $\sigma_{(1,2)}(G) = 2$ for any claw-free cubic graph G. Indeed, any adjacent pair of vertices in G is a (1,2)-spreading set in G. The results $\sigma_{(1,q)}(G) = 1$ for any $q \geq 3$ and $\sigma_{(p,q)}(G) = n(G)$ for any $p \geq 4$ and any $q \in \mathbb{N} \cup \{\infty\}$ are trivial consequences of definitions.

p	1	2	≥ 3
1	$\leq \alpha(G)(+1), n \geq 14$ [Thm. 1.3] ([Thm. 1.2])	2	1
2	u(G) + 1 or u(G) + 2 [Thm. 4.9]	u(G) or u(G) + 1 [Cor. 4.8]	u(G) or u(G) + 1 [Cor. 4.5]
3	$\beta(G)$ or $\beta(G) + 1$ [Prop. 3.12]	$\beta(G)$ [Prop. 3.8]	$\beta(G)$ [Prop. 3.8]
≥ 4	n(G)	n(G)	n(G)

The paper is organized as follows. In Section 3, we consider (3,q)-spreading numbers in claw-free cubic graphs; their values can be found in Line 3 of Table 1. (Note that 3-percolation coincides with (3,q)-spreading when $q \geq \Delta(G)$.) In Section 4, we deal with (2,q)-spreading numbers of claw-free cubic graphs; see Line 2 in Table 1. In the case of (2,q)-spreading, where $q \geq 3$, $\sigma_{(2,q)}(G) = u(G) + 1$ precisely when G is a diamond necklace N_k , while in all other claw-free cubic graphs $\sigma_{(2,q)}(G) = u(G)$. Thus, the 2-percolation number of claw-free cubic graphs is fully determined. On the other hand, when $q \leq 2$ a claw-free cubic graph can attain two values: $\sigma_{(2,2)}(G) \in \{u(G), u(G) + 1\}$, and $\sigma_{(2,1)}(G) \in \{u(G) + 1, u(G) + 2\}$, but we have not established exactly which

claw-free cubic graphs attain which of the two possible values. We conclude the paper with some open problems that arise from this study.

2. CLAW-FREE CUBIC GRAPHS

The following property of connected, claw-free, cubic graphs is established in [20].

Lemma 2.1 ([20]). If $G \neq K_4$ is a connected, claw-free, cubic graph of order n, then the vertex set V(G) can be uniquely partitioned into sets each of which induces a triangle or a diamond in G.

By Lemma 2.1, the vertex set V(G) of connected, claw-free, cubic graph $G \neq K_4$ can be uniquely partitioned into sets each of which induce a triangle or a diamond in G. Following the notation introduced in [20], we refer to such a partition as a triangle-diamond partition of G, abbreviated Δ -D-partition. We call every triangle and diamond induced by a set in our Δ -D-partition a unit of the partition. A unit that is a triangle is called a triangle-unit and a unit that is a diamond is called a diamond-unit. (We note that a triangle-unit is a triangle that does not belong to a diamond.) Two units in the Δ -D-partition are adjacent if there is an edge joining a vertex in one unit to a vertex in the other unit. In what follows, we define several structures in claw-free, cubic graphs that we will need when proving our main results. Some of these definitions have already appeared in the literature, see, for example, [4, 5, 20].

For $k \geq 2$ an integer, let N_k be the connected cubic graph constructed as follows. Take k disjoint copies D_1, D_2, \ldots, D_k of a diamond, where $V(D_i) = \{a_i, b_i, c_i, d_i\}$ and where a_ib_i is the missing edge in D_i . Let N_k be obtained from the disjoint union of these k diamonds by adding the edges $\{a_ib_{i+1} : i \in [k-1]\}$ and adding the edge a_kb_1 . We call N_k a diamond-necklace with k diamonds. Let $\mathcal{N}_{\text{cubic}} = \{N_k \mid k \geq 2\}$. The diamond-necklace, N_4 , with four diamonds is illustrated in Figure 1.

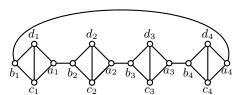


Fig. 1. The diamond-necklace N_4

For $k \geq 1$ an integer, let F_{2k} be the connected cubic graph constructed as follows. Take 2k disjoint copies T_1, T_2, \ldots, T_{2k} of a triangle, where $V(T_i) = \{x_i, y_i, z_i\}$ for $i \in [2k]$. Let

$$E_a = \{x_{2i-1}x_{2i} \colon i \in [k]\},\$$

$$E_b = \{y_{2i-1}y_{2i} \colon i \in [k]\},\$$

$$E_c = \{z_{2i}z_{2i+1} \colon i \in [k]\},\$$

where addition is taken modulo 2k (and so, $z_1 = z_{2k+1}$). Let F_{2k} be obtained from the disjoint union of these 2k triangles by adding the edges $E_a \cup E_b \cup E_c$. The resulting graph F_{2k} we call a *triangle-necklace* with 2k triangles. Let $\mathcal{T}_{\text{cubic}} = \{F_{2k} : k \geq 1\}$. The triangle-necklace F_6 is shown in Figure 2.

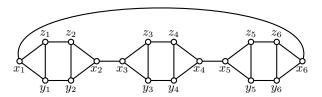


Fig. 2. The triangle-necklace F_6

For $k \geq 2$ an integer, take 2k disjoint copies T_1, T_2, \ldots, T_{2k} of a triangle, where $V(T_i) = \{x_i, y_i, z_i\}$ for $i \in [2k]$, and take k disjoint copies D_1, D_2, \ldots, D_k of a diamond, where $V(D_j) = \{a_j, b_j, c_j, d_j\}$ and where $a_j b_j$ is the missing edge in D_j for $j \in [k]$. Let

$$\begin{split} E_1 &= \{x_{2i-1}a_i \colon i \in [k]\}, \\ E_2 &= \{x_{2i}b_i \colon i \in [k]\}, \\ E_3 &= \{y_{2i-1}z_{2i+1} \colon i \in [k-1]\} \cup \{y_{2k-1}z_1\}, \\ E_4 &= \{y_{2i}z_{2i+2} \colon i \in [k-1]\} \cup \{y_{2k}z_2\}. \end{split}$$

Let H_{2k} be obtained from the disjoint union of these 2k triangles and k diamonds by adding the edges $E_1 \cup E_2 \cup E_3 \cup E_4$. The resulting graph H_{2k} we call a triangle-diamond-necklace with k diamonds. Let $\mathcal{H}_{\text{cubic}} = \{H_{2k} : k \geq 2\}$. The triangle-diamond-necklace H_6 is shown in Figure 3.

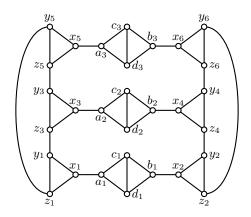


Fig. 3. The triangle-diamond-necklace H_6

3. 3-PERCOLATION IN CLAW-FREE, CUBIC GRAPHS

In this section, we study 3-percolation in claw-free, cubic graphs. The following lemma is already known in the literature and follows readily from the definition of r-bootstrap percolation.

Lemma 3.1 ([19]). For $r \geq 2$ if H is a subgraph of a graph G such that every vertex in H has strictly less than r neighbors in G that belong to $V(G) \setminus V(H)$, then every r-percolating set of G contains at least one vertex of H.

As a consequence of Lemma 3.1, we establish a relationship between the r-percolation number and the independence number of an r-regular graph. Recall that $\beta(G)$ is the cardinality of a minimum vertex cover of G.

Proposition 3.2. Let G be an r-regular graph of order n, where $r \geq 2$. A set $S \subset V(G)$ is an r-percolating set in G if and only if S is a vertex cover in G. In particular, $m(G,r) = \beta(G) = n - \alpha(G)$.

Proof. Let S be an r-percolating set of an r-regular graph G. Suppose that G-S contains an edge uv. Let $H=G[\{u,v\}]$, and so $H=K_2$ is a subgraph of G such that every vertex in H has strictly less than r neighbors in G that belong to $V(G)\setminus V(H)$. By Lemma 3.1, we infer that every r-percolating set of G contains at least one vertex of H, which is a contradiction with our assumption. Thus, S is a vertex cover of G. Conversely, suppose that S is a vertex cover of G, and let $I=V(G)\setminus S$. Since the set I is an independent set of G, every vertex in I has all its r neighbors in the set $V(G)\setminus I$. The set S is therefore an r-percolating set of G. Since m(G,r) is the cardinality of a minimum r-percolating set, we infer that $m(G,r)=\beta(G)$. Consequently, by the famous Gallai theorem $(\alpha(G)+\beta(G)=n(G))$ in any graph G0 we also have $m(G,r)=n-\alpha(G)$.

Let us recall the statement of the well-known Brooks' Coloring Theorem.

Theorem 3.3 (Brooks' Coloring Theorem). If G is a connected graph, which is not a complete graph or an odd cycle, then $\chi(G) \leq \Delta(G)$.

As a consequence of Theorem 3.3, we have the following trivial lower bound on the independence number of a graph.

Theorem 3.4. If $G \neq K_n$ is a connected graph of order n with maximum degree $\Delta \geq 3$, then $\alpha(G) \geq \frac{n}{\Delta}$.

Proof. As a consequence of Theorem 3.3, if $G \neq K_n$ is a connected graph of order n with maximum degree $\Delta \geq 3$, then the chromatic number of G is at most Δ , that is, $\chi(G) \leq \Delta$. Alternatively, viewing a Δ -coloring of G as a partitioning of its vertices into Δ independent sets, called color classes, we infer by the Pigeonhole Principle that G contains an independent set of cardinality at least n/Δ , implying that $\alpha(G) \geq n/\Delta$. \square

The following upper bound on the independence number of a claw-free graph is given by several authors.

Theorem 3.5 ([14, 23]). If G is a claw-free graph of order n with minimum degree δ , then

$$\alpha(G) \le \left(\frac{2}{\delta+2}\right)n.$$

As a consequence of Theorems 3.4 and 3.5, we have the following bounds on the independence number of a claw-free graph.

Theorem 3.6. If $G \neq K_n$ is a connected, claw-free graph of order n with minimum degree δ and maximum degree $\Delta \geq 3$, then

$$\frac{n}{\Delta} \le \alpha(G) \le \left(\frac{2}{\delta+2}\right)n.$$

As a consequence of Proposition 3.2 and Theorem 3.6, we have the following result.

Theorem 3.7. If $G \neq K_4$ is a connected, claw-free, cubic graph of order n, then the following properties hold.

- (a) $\frac{1}{3}n \le \alpha(G) \le \frac{2}{5}n$.
- (b) $\frac{3}{5}n \le m(G,3) \le \frac{2}{3}n$.

We show next that the bounds of Theorem 3.7 are tight (in the sense that they hold for connected graphs of arbitrarily large orders). Suppose that $G \in \mathcal{T}_{\text{cubic}}$. Thus, G is a triangle-necklace F_{2k} for some $k \geq 1$, and so G has order n = 6k and contains 2k vertex disjoint triangles. Moreover, every unit of the (unique) Δ -D-partition of G is a triangle-unit, implying that $\alpha(G) \leq \frac{1}{3}n$. By Theorem 3.7(a), $\alpha(G) \geq \frac{1}{3}n$. Consequently, $\alpha(G) = \frac{1}{3}n$. For example, the white vertices in the triangle-necklace F_6 of order n = 18 shown in Figure 4 form an α -set of F_6 (of cardinality $6 = \frac{1}{3}n$). By Proposition 3.2, we infer that $m(G,3) = \frac{2}{3}n$, and the corresponding set of shaded vertices in Figure 4 is a β -set of G. This shows that the lower bound of Theorem 3.7(a) is tight, as is the upper bound of Theorem 3.7(b).

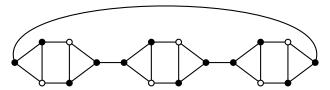


Fig. 4. A β -set in the triangle-necklace F_6

Suppose next that $G \in \mathcal{H}_{\text{cubic}}$. Thus, G is a triangle-diamond-necklace H_{2k} for some $k \geq 2$, and so G has order n = 10k. We can choose an independent set of G to contain one vertex from every triangle-unit and two vertices from every diamond-unit, implying that $\alpha(G) \geq 4k = \frac{2}{5}n$. For example, the white vertices in the triangle-diamond-necklace H_6 of order n = 30 shown in Figure 5 form an α -set of H_6 (of cardinality $12 = \frac{2}{5}n$). By Proposition 3.2, we infer that $m(G,3) = \frac{3}{5}n$, and the corresponding set of shaded vertices in Figure 5 is a β -set of H_6 . This shows that the upper bound of Theorem 3.7(a) is tight, as is the lower bound of Theorem 3.7(b).

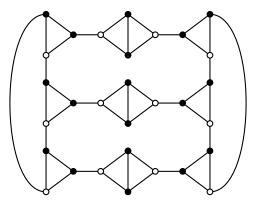


Fig. 5. A β -set in the triangle-diamond-necklace H_6

Note that m(G,3) in (claw-free) cubic graphs coincides with $\sigma_{(3,q)}(G)$ for all $q \geq 3$, establishing the (3,3)-entry of Table 1. Hence, by Proposition 3.2, we have $\sigma_{(3,q)}(G) = m(G,3) = \beta(G) = n(G) - \alpha(G)$ for all $q \geq 3$. The result can be extended to $\sigma_{(3,2)}(G)$, which yields the following result.

Proposition 3.8. If G is a connected, claw-free cubic graph of order n, then $\sigma_{(3,q)}(G) = \beta(G)$ for every $q \geq 2$.

Proof. We have already established that $\sigma_{(3,q)}(G) = \beta(G)$ for every $q \geq 3$, so it remains to resolve the case q=2. Consider a 3-percolating set P of a connected, claw-free cubic graph G, and color all its vertices blue. By Proposition 3.2, P is a vertex cover and so $V(G) \setminus P$ is an independent set in G. Now, since an independent set in G contains at most one vertex from each triangle, a vertex cover contains at least two vertices in each triangle. This implies that every vertex in P has a neighbor in P. Therefore, since $\deg_G(x)=3$ for all $x\in V(G)$, we infer that every (blue) vertex in P has at most two (white) neighbors in $V(G)\setminus P$. Thus, P is a (3,2)-spreading set. Since this is true for every 3-percolating set P of G, we infer that $\sigma_{(3,2)}(G) \leq m(G,3) = \beta(G)$.

Conversely, by Observation 1.1, $\sigma_{(3,2)}(G) \geq \sigma_{(3,3)}(G)$. Combining the obtained inequalities, we infer

$$\beta(G) \ge \sigma_{(3,2)}(G) \ge \sigma_{(3,3)}(G) = m(G,3) = \beta(G).$$

We remark that $\sigma_{(3,1)}(G) > \beta(G)$ if $G \in \mathcal{H}_{\text{cubic}}$ is a triangle-diamond necklace. Indeed, in any vertex cover P of G, every vertex in P has a neighbor in P, and so it does not have at most one neighbor in $V(G) \setminus P$.

Recall that two units in the Δ -D-partition are adjacent if there exists at least one edge joining a vertex in one unit to a vertex in the other unit. We say that a unit in the Δ -D-partition of G is *infected* if all vertices in the unit are infected.

Lemma 3.9. If G is a connected, claw-free cubic graph that contains only triangle-units, then there exists an independent set in G that contains a vertex from every triangle of G.

Proof. By supposition, every unit in the Δ -D-partition of the connected, claw-free cubic G is a triangle-unit. Since G is diamond-free, we note that the triangle-units in G correspond to the triangles in G. The graph G has order n=3t where t denotes the number of triangle-units in G. By Theorem 3.7(a), we have $t=\frac{1}{3}n\leq\alpha(G)$. Since every independent set in G contains at most one vertex from every triangle-unit in G, we infer that $\alpha(G)\leq t=\frac{1}{3}n$. Consequently, $\alpha(G)=\frac{1}{3}n=t$ and every maximum independent set in G contains a vertex from every triangle of G.

Lemma 3.10. If G is a connected, claw-free cubic graph, then there exists an independent set in G that contains a vertex from every triangle of G.

Proof. If the connected, claw-free cubic G that contains only triangle-units, then the result follows from Lemma 3.9. Hence, we may assume that G contains at least one diamond-unit. We now construct an independent set I of G as follows. Initially, we set $I = \emptyset$. For each diamond-unit D of G, we add to the set I exactly one of the two vertices of degree 3 in the diamond-unit D. Thereafter, we delete all diamond-units from the graph G. If G contains only diamond-units, then the resulting set I has the desired property that it contains a vertex from every triangle of G. Hence, we may assume that G contains at least one triangle-unit.

Let G' be the graph obtained by deleting all diamond-units from G. We note that every component, C, of G_1 has the following properties: (1) the component C contains no diamond, (2) every vertex in C belongs to a triangle in the component C and (3) the component C has minimum degree 2 (and maximum degree at most 3). We now select an arbitrary vertex u_1 of degree 2 in C, add the vertex u_1 to the set I, and delete the triangle that contains the vertex u_1 . In the resulting graph G_2 , once again properties (1), (2) and (3) hold, and we select an arbitrary vertex u_2 of degree 2 in G_2 , add the vertex u_2 to the set I, and delete from G_2 the triangle that contains the vertex u_2 . Upon completion of this process, the resulting set I is an independent set in G that contains a vertex from every triangle of G.

We note that the statement in Lemma 3.10 is equivalent to the following statement.

Lemma 3.11. If G is a connected, claw-free cubic graph, then there exists a vertex cover P such that every triangle of G contains exactly two vertices from P.

Proposition 3.12. If G is a connected, claw-free cubic graph, then

$$\sigma_{(3,1)}(G) \le \beta(G) + 1,$$

and this bound is sharp.

Proof. Let G be a connected, claw-free cubic graph, and let P be a 3-percolating set of G satisfying $|P| = m(G,3) = \beta(G)$. We note that P is a vertex cover of G and $V(G) \setminus P$ is a maximum independent set of G. Clearly, every vertex, which is not in P, has three neighbors in P, and so the first condition of the (3,1)-spreading rule is always satisfied with respect to P. In addition, due to Lemma 3.11 we may assume that P contains exactly two vertices from every triangle in G, or, equivalently, so that the independent set $V(G) \setminus P$ contains a vertex from every triangle in G. Now,

consider an arbitrary triangle T in G, and let $P^* = P \cup \{v\}$, where v is the unique vertex in $V(T) \setminus P$. Let the vertices of P^* be initially infected, and note that (with respect to P^*) the triangle T is completely infected.

Let u be a vertex that is not yet infected and is adjacent to an infected triangle T' (with all three vertices of T' infected). In particular, we note that $u \notin P^*$. Suppose firstly that u is adjacent to two vertices of T'. Thus, the vertices in $V(T') \cup \{u\}$ induce a diamond-unit. In this case, the vertex u immediately becomes infected, since both of its neighbors in T' have no other uninfected neighbor. As a result, the (unique) triangle containing the vertex u becomes an infected triangle, and so the diamond-unit containing u is infected. Suppose next that u is adjacent to exactly one vertex u' of T'. Once again in this case, the vertex u immediately becomes infected, since its neighbor u' in u' has no other uninfected neighbor. On the other hand, suppose that u' is infected and is adjacent to an infected triangle u' (with all three vertices of u' infected), and there exists an uninfected neighbor u' of u' that belongs to the same unit as u'. Then u' becomes infected, because u' has only one uninfected neighbor, namely u'.

In all of the above cases, the unit, which is adjacent to an infected unit, becomes infected. Since G is connected, and initially there is a triangle infected (making the unit in which the triangle lies also infected), we infer that due to the above arguments all vertices in G become infected. Thus, $\sigma_{(3,1)}(G) \leq |P^*| = |P| + 1 = \beta(G) + 1$, thereby proving the desired upper bound. The case of a triangle-diamond necklace $G \in \mathcal{H}_{\text{cubic}}$ shows that the bound is sharp.

4. 2-PERCOLATION IN CLAW-FREE, CUBIC GRAPHS

In this section, we study 2-percolation in claw-free, cubic graphs. It is easy to see that $m(K_4, 2) = 2$. In what follows we therefore restrict our attention to 2-percolation in connected, claw-free, cubic graphs G different from K_4 . Thus, by Lemma 2.1, such a graph G has a Δ -D-partition (in which every unit is a triangle-unit or a diamond-unit). Recall that u(G) is defined as the number of units in this (unique) Δ -D-partition. In what follows, if D is a diamond-unit in the Δ -D-partition of G, then a dominating vertex in D is a vertex that is adjacent to the three other vertices in the diamond-unit D.

Proposition 4.1. If
$$G \in \mathcal{N}_{\text{cubic}}$$
, then $m(G,2) = u(G) + 1$.

Proof. Suppose that $G \in \mathcal{N}_{\text{cubic}}$, and so G is a diamond-necklace N_k for some $k \geq 2$. Thus, G has u(G) = k units, each of which is a diamond-unit. Let S be a 2-percolating set of G of minimum cardinality, and so S is an 2-percolating set of G satisfying |S| = m(G, 2). By Lemma 3.1, we infer that the set S contains at least one vertex from every diamond-unit of G. Moreover, if D is a diamond-unit of G and the set S contains exactly one vertex from D, then by Lemma 3.1 we infer that such a vertex of S is a dominating vertex of D. In particular, since S contains at least one vertex from every unit in the Δ -D-partition of G, we note that $m(G, 2) = |S| \geq u(G) = k$.

Suppose that |S| = k. By our earlier observations, the set S contains exactly one vertex from every diamond-unit, namely a vertex from each diamond-unit that dominates that unit. The resulting set S is a 2-packing in G; that is, $d_G(u, v) \geq 3$ for

every two distinct vertices u and v that belong to S. Moreover in this case, letting H = G - S, we note that H is a cycle C_{3k} and every vertex in H has exactly one neighbor that belongs to $V(G) \setminus V(H) = S$. Thus by Lemma 3.1, we infer that every 2-percolating set of G contains at least one vertex of H. However, this contradicts our supposition that S is a 2-percolating set of G that contains no vertex of G. Hence, G and G is a 2-percolating set of G that contains no vertex of G.

To establish an upper bound on m(G,2) in this case when $G=N_k$, if S^* consists of a dominating vertex from k-1 diamond-units and two dominating vertices from the remaining diamond-unit, then S^* is a 2-percolating set of G and $|S^*|=k+1$. For example, if $G=N_4$ (here k=4), then such a set S^* illustrated in Figure 6 by the shaded vertices is a 2-percolating set of G and $|S^*|=5$.

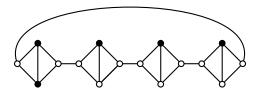


Fig. 6. The diamond-necklace N_4 and a 2-percolating set of size 5

This implies that $m(G,2) \leq |S^*| = k+1$. Consequently, m(G,2) = k+1. Thus in this case when $G \in \mathcal{N}_{\text{cubic}}$, we have shown that m(G,2) = u(G) + 1.

If two distinct units are joined by at least two edges, then we say that these two units are double-bonded.

Lemma 4.2. If $G \neq K_4$ is a connected, claw-free, cubic graph and $G \notin \mathcal{N}_{\text{cubic}}$, then there exists a triangle-unit T', such that the graph G - T' has at most two components.

Proof. Let \mathcal{T} be the set of all triangle-units in G. Suppose to the contrary that for every $T \in \mathcal{T}$ the graph G - T consists of three components. Let $T_1 \in \mathcal{T}$ be a triangle-unit such that at least one of the three components of the graph $G - T_1$ does not contain a triangle-unit. Denote this component by G_1 and let D_1 be the diamond-unit in G_1 adjacent to T_1 . Then in G_1 , D_1 is adjacent to exactly one diamond-unit D_2 , which is in turn adjacent to another diamond-unit D_3 and so on, which yields a contradiction since G is finite. Therefore, G_1 contains a triangle-unit $T' \notin \mathcal{T}$, which is the final contradiction.

The following observation will be necessary for proving the subsequent theorem.

Observation 4.3. Let $G \neq K_4$ be a connected, claw-free, cubic graph, and U, V adjacent units in the Δ -D-partition of G. Also, let uv be an edge connecting units U and V, where $u \in U$ and $v \in V$. If u is infected and another vertex $v_1 \in V$, where $d(v_1, u) = 2$, is also infected, then the whole unit V becomes infected.

Proof. Vertex v has two infected neighbors, namely u and v_1 . After that, if V is a triangle unit, the remaining vertex of V is also infected, and if V is a diamond unit, the remaining two vertices of V become infected after two steps.

Theorem 4.4. If $G \neq K_4$ is a connected, claw-free, cubic graph of order n that contains u(G) units, then

$$m(G,2) = \begin{cases} u(G) + 1, & \text{if } G \in \mathcal{N}_{\text{cubic}}, \\ u(G), & \text{otherwise}. \end{cases}$$

Proof. By Lemma 3.1, we infer that every 2-percolating set of G contains at least one vertex from every triangle-unit and every diamond-unit of G. Moreover, if D is a diamond-unit of G and a 2-percolating set contains exactly one vertex from D, then by Lemma 3.1 we infer that such a vertex is a dominating vertex of D. In particular, since every 2-percolating set of G contains at least one vertex from every unit in the Δ -D-partition of G, we note that $m(G,2) \geq u(G)$. If $G \in \mathcal{N}_{\text{cubic}}$, then by Proposition 4.1, m(G,2) = u(G) + 1. Hence, we may assume that $G \notin \mathcal{N}_{\text{cubic}}$, for otherwise the desired result follows.

Since every triangle-unit contributes 3 to the order of the graph and every diamond-unit contributes 4 to the order of the graph, we observe that if G has order n with u_t triangle-units and u_d diamond-units, then $u(G) = u_t + u_d$ and $n = 3u_t + 4u_d$. By assumption, $u_t \geq 1$. Thus, since n is even, $u_t \geq 2$, that is, G contains at least two triangle-units. According to Lemma 4.2, there exists a triangle-unit T_1 of G, where $V(T_1) = \{t_1, t_2, t_3\}$, such that $G - T_1$ consists of at most two components. This implies that at most one of the edges incident with exactly one vertex of T_1 is a bridge. Since G is connected, the triangle-unit T_1 is adjacent to at least one other unit. Let U_1 be a unit that is adjacent to T_1 such that the edge between T_1 and T_2 is not a bridge, or T_1 and T_2 are double-bonded. Finally denote as T_2 the component of T_2 containing T_2 .

As observed earlier, $m(G,2) \geq u(G)$. Hence, it suffices for us to show that $m(G,2) \leq u(G)$. For this purpose, we construct a 2-percolating set S of G that contains exactly one vertex from each unit of G. Initially, we let $S = \emptyset$. We consider three cases. First, suppose that the units T_1 and T_2 are not double-bonded.

Case 1. The unit U_1 is a triangle-unit. Let $V(U_1) = \{a_1, b_1, c_1\}$ and where t_1c_1 is an edge. In this case, we add to S the vertices t_1 and a_1 , and so $S = \{t_1, a_1\}$. The vertices in S are indicated by the shaded vertices in Figure 7(a). Due to Observation 4.3, every vertex in the triangle-unit U_1 becomes infected.

Case 2. The unit U_1 is a diamond-unit. Let $V(U_1) = \{a_1, b_1, c_1, d_1\}$ and where a_1b_1 is the missing edge in U_1 and where t_1b_1 is an edge of G. In this case, we add to S the vertices t_1 and c_1 , and so $S = \{t_1, c_1\}$. The vertices in S are indicated by the shaded vertices in Figure 7(b). Due to Observation 4.3, every vertex in the diamond-unit U_1 becomes infected.

Case 3. The units T_1 and U_1 are double-bonded. Thus, there is a vertex $u \in U_1$ that is joined to a vertex of T_1 different from t_1 . Renaming the vertices t_2 and t_3 if necessary, we may assume that ut_2 is such an edge between the units T_1 and U_1 . In particular,

we note that $t_2 \notin S$ and let $S = \{t_1, u\}$. The vertices in S are indicated by the shaded vertices in Figure 7(c). Again, due to Observation 4.3, every vertex in both units T_1 and U_1 becomes infected. In this case, the set S consists of two vertices, one vertex from each of the units U_1 and T_1 .

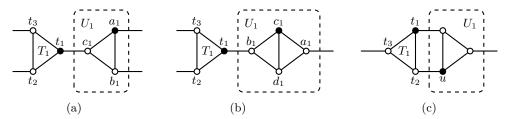


Fig. 7. Possible adjacent units in the proof of Theorem 4.4

Cases 1, 2 and 3 occur in Step 1 of our procedure to construct the set S. Thus after Step 1, the initial set S consists of two vertices, namely a vertex t_1 from the triangle-unit T_1 and one vertex from the unit U_1 that is adjacent to T_1 . Moreover in all three cases, every vertex in the unit U_1 becomes infected. We now proceed to Step 2 of our procedure to construct the set S.

Let U_2 be a unit different from T_1 that is adjacent to the infected unit U_1 , and let uv be an edge that joins a vertex $u \in U_1$ and a vertex $v \in U_2$. Note that if there is no such unit U_2 , then G consist only of units T_1 and U_1 and is now fully infected, yielding the desired result. Since $U_2 \neq T_1$, we note that $v \neq t_1$. In particular, we note that $v \notin S$ since at this stage of the construction the set S only contains the vertex t_1 and one other vertex, namely a vertex from U_1 .

Suppose that U_2 is a triangle-unit. Let $V(U_2) = \{v, v_1, v_2\}$. We now add to the set S exactly one of the vertices v_1 and v_2 . By symmetry, we may assume that v_1 is added to the set S. Once again, due to Observation 4.3, every vertex in the triangle-unit U_2 becomes infected.

Suppose next that U_2 is a diamond-unit. Let $V(U_2) = \{v, v_1, v_2, v_3\}$, where v_1 and v_2 are the dominating vertices of U_2 (and so, vv_3 is the missing edge in D_v). We now add to the set S exactly one of the dominating vertices of U_2 . By symmetry, we may assume that dominating vertex v_1 of U_2 is added to S. We again invoke Observation 4.3, thus every vertex in the diamond-unit U_2 becomes infected.

After Step 2 of our procedure to construct the set S, every vertex in the unit U_2 becomes infected, and the set S contains exactly one vertex from each of the units T_1, U_1 and U_2 .

Suppose that $i \geq 2$ and after Step i of our procedure to construct the set S, the units U_1, U_2, \ldots, U_i are all infected, and the set S contains exactly one vertex from each of the units U_1, U_2, \ldots, U_i . Moreover, if $T_1 = U_j$ for some $j \in [i] \setminus \{1\}$, then after Step i we have |S| = i and the set S contains exactly one vertex from each of the units U_1, U_2, \ldots, U_i , while if $T_1 \neq U_j$ for any $j \in [i]$, then |S| = i + 1 and S contains one vertex from each of the units $U_1, U_2, \ldots, U_i, T_1$.

In Step i+1, let U_{i+1} be a unit different the units U_1, U_2, \ldots, U_i that is adjacent to an infected unit U_j for some $j \in [i]$, and let uv be an edge that joins a vertex $u \in U_j$ and a vertex $v \in U_{i+1}$.

Suppose that U_{i+1} is a triangle-unit. Let $V(U_{i+1}) = \{v, v_1, v_2\}$. In this case, either $U_{i+1} = T_1$ and renaming vertices of T_1 if necessary we may assume that $v_1 = t_1$, or $U_{i+1} \neq T_1$, and we add to the set S the vertex v_1 . Suppose next that U_{i+1} is a diamond-unit. Let $V(U_{i+1}) = \{v, v_1, v_2, v_3\}$, where v_1 and v_2 are the dominating vertices of U_{i+1} (and so, vv_3 is the missing edge in D_v). We now add to the set S the dominating vertex v_1 of U_{i+1} . In either of these cases the vertices u and v_1 satisfy the conditions of Observation 4.3, therefore the whole unit U_{i+1} becomes infected.

Continuing in the described way, the above process continues until every unit of G_1 and also the unit T_1 become infected. If $G-T_1$ is connected, then all vertices of G have become infected and S contains exactly one vertex from each unit of G, thus the proof is complete. Otherwise, if $G-T_1$ is not connected, then it has at most two components, and let G_2 be the component of $G-T_1$ different from G_1 . Let U'_1 be the unit in G_2 that is adjacent to T_1 , and let v_2a be the bridge of G connecting T_1 with U'_1 . Now, adding a dominating vertex b of U'_1 , different from a, we infer that all vertices of U'_1 become infected (using the analogous arguments as in the previous cases). Now, we can continue with the same process, continuing by infecting the units U'_i in G_2 that are adjacent to an already infected unit by adding to S appropriately selected dominating vertex of U'_i . Since G_2 is connected, we infer that all vertices of G_2 become infected, where we put in S exactly one vertex of each unit. Thus, S is a 2-percolating set of G, implying that $m(G,2) \leq u(G)$. As observed earlier, $m(G,2) \geq u(G)$. Consequently, m(G,2) = u(G). This completes the proof of Theorem 4.4.

By definition of a (p,q)-spreading set in a graph, if G is cubic graph, then a set is a (2,3)-spreading set if and only if it is 2-percolating set, implying that $\sigma_{(2,3)}(G) = m(G,2)$. As an immediate consequence of Theorem 4.4, we infer the following result.

Corollary 4.5. If $G \neq K_4$ is a connected, claw-free, cubic graph of order n that contains u(G) units, then

$$\sigma_{(2,3)}(G) = \begin{cases} u(G) + 1, & \text{if } G \in \mathcal{N}_{\text{cubic}}, \\ u(G), & \text{otherwise.} \end{cases}$$

We are now in a position to establish the following relationships between the (2,2)-spreading number and the (2,3)-spreading number of a connected, claw-free, cubic graph.

Proposition 4.6. If $G = K_4$ or if $G \in \mathcal{N}_{\text{cubic}}$, then $\sigma_{(2,2)}(G) = \sigma_{(2,3)}(G)$.

Proof. If $G = K_4$, then $\sigma_{(2,2)}(G) = \sigma_{(2,3)}(G) = 2$. Hence, we may assume that $G \neq K_4$. Suppose that $G \in \mathcal{N}_{\text{cubic}}$. By Corollary 4.5, we have $\sigma_{(2,3)}(G) = u(G) + 1$. Moreover, as shown in the proof of Proposition 4.1, if $G = N_k$ for some $k \geq 2$, then we can choose the (2,3)-spreading set S of G to consists of a dominating vertex from k-1

diamond-units and two dominating vertices from the remaining diamond-unit in G (as illustrated in Figure 6 in the case when k=4). However, such a set S is also a (2,2)-spreading set S of G, and so $\sigma_{(2,2)}(G) \leq |S| = \sigma_{(2,3)}(G)$. By Observation 1.1, $\sigma_{(2,3)}(G) \leq \sigma_{(2,2)}(G)$. Consequently, if $G \in \mathcal{N}_{\text{cubic}}$, then $\sigma_{(2,2)}(G) = \sigma_{(2,3)}(G)$.

Proposition 4.7. If G is a connected, claw-free, cubic graph, then

$$\sigma_{(2,2)}(G) \le \sigma_{(2,3)}(G) + 1.$$

Proof. Let G be a connected, claw-free, cubic graph. If $G = K_4$ or if $G \in \mathcal{N}_{\text{cubic}}$, then by Proposition 4.6 we have $\sigma_{(2,2)}(G) = \sigma_{(2,3)}(G)$. Hence, we may assume that $G \neq K_4$ and $G \notin \mathcal{N}_{\text{cubic}}$. To prove that $\sigma_{(2,2)}(G) \leq \sigma_{(2,3)}(G) + 1$, let S be a (2,3)-spreading set of G as constructed in the proof of Theorem 4.4. Adopting the notation from the proof of Theorem 4.4, in Step 1 of our procedure to construct the set S we add one additional vertex to the set S as follows. If the unit U_1 is a triangle-unit (see Figure 7(a)), then we add the common neighbor of t_1 and a_1 , namely the vertex c_1 , to the set S. If the unit U_1 is a diamond-unit (see Figure 7(b)), then we add the common neighbor of t_1 and t_2 , namely the vertex t_2 , to the set t_2 . Let t_2 0 denote the resulting set upon completion of the construction of the set t_2 0, and so t_2 1 in Case 1 or the vertex t_2 2, the set t_2 3 is a t_2 4. With the addition of the vertex t_2 6 in Case 1 or the vertex t_2 7 in Case 2, the set t_2 7 is a t_2 7 spreading set t_2 8 of t_2 9, implying that t_2 9, where t_2 9 is a t_2 9, implying that t_2 9, where t_2 9 is a t_2 9 in Case 2, the set t_2 9 is a t_2 9 spreading set t_2 9 of t_2 9 in Case 1 or the vertex t_2 9 in Case 2, the set t_2 9 is a t_2 9 spreading set t_2 9 of t_2 9 in Case 1 or the vertex t_2 9 in Case 2, the set t_2 9 is a t_2 9 spreading set t_2 9 of t_2 9 in Case 1 or the vertex t_2 9 in Case 2, the set t_2 9 is a t_2 9 spreading set t_2 9 in Case 1 or the vertex t_2 9 in Case 2, the set t_2 9 is a t_2 9 spreading set t_2 9 in Case 3 in Case

Every (2,3)-spreading set of G is by definition also a (2,2)-spreading set of G, and so $\sigma_{(2,3)}(G) \leq \sigma_{(2,2)}(G)$. Hence, as a consequence of Proposition 4.7 and applying also Corollary 4.5, we have the following result.

Corollary 4.8. If G is a connected, claw-free, cubic graph, then

$$\sigma_{(2,3)}(G) \le \sigma_{(2,2)}(G) \le \sigma_{(2,3)}(G) + 1.$$

In addition,

$$u(G) \le \sigma_{(2,2)}(G) \le u(G) + 1.$$

If $G = F_{2k} \in \mathcal{T}_{\text{cubic}}$ then it is not difficult to see that $S = \{z_i, : i \in [2k-1]\} \cup \{x_{2k}\}$ is a (2,2)-spreading set. Notably, the process starts with the infection of vertex x_1 , and then continues with y_1, y_2 and so on. Therefore, $\sigma_{(2,2)}(G) = u(G)$. On the other hand, if $G \in \mathcal{N}_{\text{cubic}}$, then $\sigma_{(2,2)}(G) = u(G) + 1$. In both of these cases it holds that $\sigma_{(2,2)}(G) = \sigma_{(2,3)}(G)$, however this does not hold in general. For instance, see the graph G in Figure 8 for which we have $\sigma_{(2,3)} = u(G)$, where the shaded vertices indicated in the figure form a (2,3)-spreading set of G. Yet, one can verify that $\sigma_{(2,2)} = u(G) + 1$.

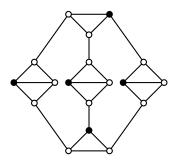


Fig. 8. A graph G with $\sigma(2,3) = u(G)$ and $\sigma(2,2) = u(G) + 1$

Finally, we consider (2,1)-spreading, and bound the corresponding invariant in claw-free cubic graphs. Again we see that it can achieve only two possible values.

Theorem 4.9. If $G \neq K_4$ is a connected, claw-free, cubic graph of order n that contains u(G) units, then

$$u(G) + 1 \le \sigma_{(2,1)}(G) \le u(G) + 2.$$

Proof. First note that due to Lemma 3.1 a (2,1)-spreading set S needs to contain at least one vertex in each unit of $G \neq K_4$. Now, if S is a set of vertices that contains exactly one vertex from each unit of G, then every vertex $v \in S$ has at least two neighbors in $V(G) \setminus S$, and therefore such a set S cannot be a (2,1)-spreading set. Hence, $u(G) + 1 \leq \sigma_{(2,1)}(G)$.

We show next that $\sigma_{(2,1)}(G) \leq u(G) + 2$ by constructing a (2,1)-spreading set S of G satisfying $|S| \leq u(G) + 2$. Let T be an arbitrary triangle in G (that may also be part of a diamond-unit). Initially, we let S = V(T), and so the vertices in T form the set of initially infected vertices. We now enlarge the set S as follows. Suppose that T belongs to a diamond-unit D. In this case, we let v be the vertex in the diamond-unit D that is not in T. The vertex v immediately becomes infected in the (2,1)-spreading process, since both of its neighbors in T have no other uninfected neighbor. As a result, all vertices in the diamond-unit D become infected, that is, D becomes an infected unit. Suppose next that T is a triangle-unit of G. Then T is already an infected unit.

Finally let T' be the infected unit containing T (that is, T' = D or T' = T). We now extend the set S in the same manner as in the proof of Theorem 4.4. Let U be a unit adjacent to T'. If U is a triangle-unit, then we add to S a vertex in U that is not adjacent to any vertex in the unit T'. If U is a diamond-unit, then we add to S a vertex in U that is a dominating vertex of the unit U. Proceeding analogously as in the proof of Theorem 4.4, this results in all vertices of the unit U becoming infected in the (2,1)-spreading process.

Continuing in this way by considering a vertex not yet infected that is adjacent to an infected triangle, we obtain a (2,1)-spreading set S of G starting with the set V(T) and adding exactly one vertex from every unit of G that does not contain the triangle T. Thus, $\sigma_{(2,1)}(G) \leq |S| = u(G) + 2$.

5. CONCLUDING REMARKS

We end this paper with some remarks concerning the computational complexity of determining $\sigma_{(p,q)}(G)$ in claw-free cubic graphs G. We have already mentioned that it is unclear whether one can determine the zero forcing number of a claw-free cubic graph efficiently. For some of the values of p and q, where p > 1 or q > 1, there exists a polynomial-time algorithm to determine $\sigma_{(p,q)}(G)$. First, determining the independence number in claw-free graphs can be done in polynomial time (see [24], where a polynomial-time algorithm is presented even for the weighted version of the problem). Therefore, the problem of determining $\sigma_{(3,q)}(G)$ is polynomial in claw-free cubic graphs G for all $q \geq 2$. Similarly, one can efficiently determine the number of units in a claw-free cubic graph. Thus, the problem of determining $\sigma_{(2,q)}(G)$ is also polynomial in claw-free cubic graphs G for any $q \geq 3$.

The following problems remain open:

Problem 5.1. Provide a structural characterization of the connected claw-free cubic graphs G for which $\sigma_{(2,2)} = u(G)$ (resp., u(G)+1). Is there a polynomial-time algorithm to recognize the corresponding classes of connected claw-free cubic graphs?

Problem 5.2. Provide a structural characterization of the connected claw-free cubic graphs G for which $\sigma_{(2,1)} = u(G) + 1$ (resp., u(G) + 2). Is there a polynomial-time algorithm to recognize the corresponding classes of connected claw-free cubic graphs?

Problem 5.3. Provide a structural characterization of the connected claw-free cubic graphs G for which $\sigma_{(3,1)} = \beta(G)$ (resp., $\beta(G) + 1$). Is there a polynomial-time algorithm to recognize the corresponding classes of connected claw-free cubic graphs?

Acknowledgements

We are grateful to anonymous reviewer for meticulous reading of the first version of the manuscript and for numerous useful suggestions. Boštjan Brešar was supported by the Slovenian Research and Innovation Agency (ARIS) under the grants P1-0297, N1-0285, and J1-4008. Jaka Hedžet was supported by the Slovenian Research and Innovation Agency (ARIS) under the grant P1-0297. Michael A. Henning was supported in part by the South African National Research Foundation (grants 132588, 129265) and the University of Johannesburg.

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Received: December 18, 2024. Revised: July 6, 2025.

Accepted: July 11, 2025.