

ON THE LONGEST PATH IN A RECURSIVELY PARTITIONABLE GRAPH

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Abstract. A connected graph G with order $n \geq 1$ is said to be recursively arbitrarily partitionable (R-AP for short) if either it is isomorphic to K_1 , or for every sequence (n_1, \dots, n_p) of positive integers summing up to n there exists a partition (V_1, \dots, V_p) of $V(G)$ such that each V_i induces a connected R-AP subgraph of G on n_i vertices. Since previous investigations, it is believed that a R-AP graph should be “almost traceable” somehow. We first show that the longest path of a R-AP graph on n vertices is not constantly lower than n for every n . This is done by exhibiting a graph family \mathcal{C} such that, for every positive constant $c \geq 1$, there is a R-AP graph in \mathcal{C} that has arbitrary order n and whose longest path has order $n - c$. We then investigate the largest positive constant $c' < 1$ such that every R-AP graph on n vertices has its longest path passing through $n \cdot c'$ vertices. In particular, we show that $c' \leq \frac{2}{3}$. This result holds for R-AP graphs with arbitrary connectivity.

Keywords: recursively partitionable graph, longest path.

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1. INTRODUCTION

Let $n \geq 1$ be a positive integer. A n -graph is a graph whose *order*, i.e. its number of vertices, is n . Throughout this paper, we denote by $LP(G)$ the order of the longest path in a given connected graph G . We say that G is *recursively arbitrarily partitionable* (R-AP for short) if and only if one of the following two conditions hold.

- The graph G is an isolated vertex.
- For every sequence (n_1, \dots, n_p) of positive integers that performs a partition of n , there exists a partition (V_1, \dots, V_p) of $V(G)$ such that $G[V_i]$ is a connected R-AP subgraph of G on n_i vertices for all $i \in \{1, \dots, p\}$.

The property of being R-AP was introduced in [7] as a strengthened version of the property of being *arbitrarily partitionable*. The property of being AP was itself

introduced to deal with a problem of resource sharing among an arbitrary number of users (see [1, 2, 5, 8] for further details).

R-AP graphs have been mainly studied in the context of some simple classes of graphs like trees [7], a family of unicyclic 1-connected graphs called *suns* [6], and a class of 2-connected graphs called *balloons* [4, 7]. Although these works did not lead to numerous general properties of R-AP graphs, they however suggest that the property of being R-AP is even closer to *traceability*¹⁾ than the one of being AP. For instance, we know that if T is a R-AP n -tree, then $LP(T) \geq n - 2$. It was also empirically observed²⁾ that if B is a R-AP n -balloon, then $LP(B) \geq n - 4$. Such bounds do not exist regarding AP trees and AP balloons since the structure of these graphs is much less predictable (see [3] and [4], respectively).

Regarding these observations, one could naively think that there should exist a small positive constant $c \geq 1$ such that $LP(G) \geq n - c$ for every R-AP n -graph G . In this work, we first show, in Section 3, that such a constant does not exist by exhibiting a class \mathcal{C} of R-AP graphs such that for every c there exists a n -graph C in \mathcal{C} such that $LP(C) = n - c$ for some n . The graphs of \mathcal{C} are 1-connected, but an equivalent result regarding 2-connected graphs is derived by slightly modifying our construction. We then investigate, in concluding Section 4, the greatest constant $c' \leq 1$ such that every R-AP n -graph has its longest path passing through $n \cdot c'$ of its vertices. In particular, we exhibit another family of graphs showing that $c' \leq \frac{2}{3}$. This upper bound also holds regarding ℓ -connected R-AP graphs, no matter what is the value of ℓ .

2. DEFINITIONS AND PRELIMINARY RESULTS

First observe that adding edges to a R-AP graph does not make it loose its property of being R-AP.

Remark 2.1. If G is spanned by a R-AP subgraph, then G is R-AP.

Because every path is clearly R-AP, the next result follows by Remark 2.1.

Remark 2.2. Every traceable graph is R-AP.

Determining whether a n -graph G is R-AP is laborious since, according to the original definition, one has to check whether G can be partitioned following every partition of n . We thus usually prefer to check the following equivalent condition which derives from the fact that a R-AP graph is partitionable into R-AP subgraphs at will.

Remark 2.3 ([7]). A connected n -graph G is R-AP if and only if for every $\lambda \in \{1, \dots, \lfloor \frac{n}{2} \rfloor\}$ there exists a partition $(V_\lambda, V_{n-\lambda})$ of $V(G)$ such that $G[V_\lambda]$ and $G[V_{n-\lambda}]$ are connected R-AP subgraphs of G on λ and $n - \lambda$ vertices, respectively.

Let us now introduce the following subclass of *caterpillar graphs*.

¹⁾ A *traceable* graph is a graph that has a Hamiltonian path.

²⁾ Private communication.

Definition 2.4. Let $a, b \geq 2$ be two positive integers and consider three vertex-disjoint paths u_1u_2, v_1, \dots, v_a and w_1, \dots, w_b of order 2, a and b , respectively. The *caterpillar* $Cat(a, b)$ is the tree obtained by identifying the vertices u_1, v_1 and w_1 .

Throughout this paper, every mention to caterpillar graphs actually refers to caterpillars of the form $Cat(a, b)$. Two examples of such caterpillars are given in Figure 1. This family of caterpillars is important regarding R-AP graphs since it was proven in [7] that most of R-AP trees are caterpillars of this kind. The authors of [7] also gave a complete characterization of R-AP caterpillars.



Fig. 1. The caterpillars $Cat(2, 3)$ and $Cat(3, 3)$

Theorem 2.5 ([7]). *A caterpillar $Cat(a, b)$ is R-AP if and only if a and b take values in Table 1.*

Table 1. Values a and b ($a \leq b$) such that $Cat(a, b)$ is R-AP

a	b
2, 4	$\equiv 1 \pmod 2$
3	$\equiv 1, 2 \pmod 3$
5	6, 7, 9, 11, 14, 19
6	7
7	8, 9, 11, 13, 15

3. LONGEST PATH AND ADDITIVE FACTOR

In this section, we prove the following result.

Theorem 3.1. *There does not exist a positive constant $c \geq 1$ such that we have $LP(G) \geq n - c$ for every R-AP n -graph G .*

This is proved by exhibiting a counterexample for every possible value of c . For this purpose, we introduce the family of *connected cycles* graphs.

Definition 3.2. Let $k \geq 1$ and $x, y \geq 0$ be three positive integers. The *connected cycles* graph $CC_k(x, y)$ is the graph with the following vertices:

- Let $u_1 \dots u_x$ and $v_1 \dots v_y$ be paths with order x and y , respectively.
- For every $i \in \{1, \dots, k\}$, let $a_i b_i e_i d_i c_i a_i$ be a cycle with length 5.
- For every $i \in \{1, \dots, k - 1\}$, let $w_{i, i+1}$ be a vertex.

These vertices are linked in $CC_k(x, y)$ in the following way: $u_x a_1, v_y e_k \in E(CC_k(x, y))$ and we have $w_{i,i+1} e_i, w_{i,i+1} a_{i+1} \in E(CC_k(x, y))$ for every $i \in \{1, \dots, k - 1\}$.

An example of a connected cycles graph is depicted in Figure 2. Notice that $LP(CC_k(1, 1)) = |V(CC_k(1, 1))| - k$. Thus, by showing that all graphs $CC_k(1, 1)$ are R-AP, we can contradict the existence of the constant c mentioned in Theorem 3.1.

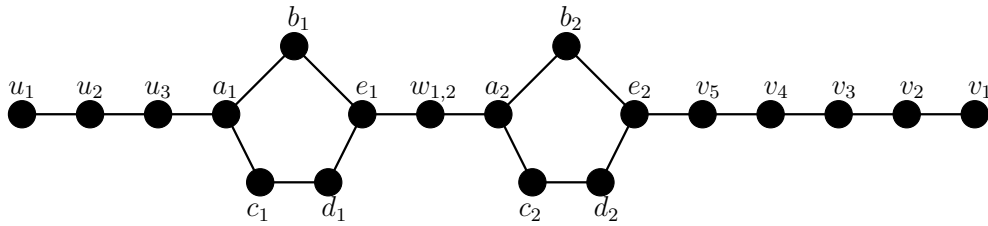


Fig. 2. The connected cycles graph $CC_2(3, 5)$

Before proving that $CC_k(1, 1)$ is R-AP for every k , we first introduce another graph structure we encounter while partitioning a connected cycles graph.

Definition 3.3. Let $k \geq 1$ and $x \geq 0$ be two positive integers. The *partial connected cycles* graph $PCC_k(x)$ is the graph obtained by removing the vertex e_k from $CC_k(x, 0)$.

We are now ready to prove the main result of this section.

Lemma 3.4. *The graph $PCC_k(x)$ is R-AP for every $k \geq 1$ and $x \geq 1$ such that $x \not\equiv 2 \pmod 3$. The graph $CC_k(x, y)$ is R-AP for every $k \geq 1$ and $x, y \geq 1$ whenever $x \not\equiv 2 \pmod 3$ or $y \not\equiv 2 \pmod 3$.*

Proof. The proof is by induction on k and uses the terminology introduced in Definition 3.2. For each value of k , we prove that the result is true for all possible values of x and (possibly) y which satisfy the claim. Recall that, according to Remark 2.3, a connected n -graph G is R-AP if and only if for every $\lambda \in \{1, \dots, \lfloor \frac{n}{2} \rfloor\}$ we can partition $V(G)$ into two parts V_λ and $V_{n-\lambda}$ inducing connected R-AP subgraphs of G with order λ and $n - \lambda$, respectively.

Case 1. $k = 1$.

First, every graph $PCC_1(x)$ is R-AP since it is spanned by $Cat(3, x + 1)$, which is R-AP according to the assumption on x .

We now prove that every graph $C = CC_1(x, y)$ is R-AP whenever the conditions of the claim are fulfilled. This is proved by induction on $x + y$ by showing that there is a partition of $V(C)$ into two parts V_λ and $V_{n-\lambda}$ satisfying the conditions above for every $\lambda \in \{1, \dots, \lfloor \frac{n}{2} \rfloor\}$ where $n = 5 + x + y$. For each value of λ , we give a satisfying subset V_λ , and it is understood that $V_{n-\lambda} = V(C) - V_\lambda$. We further assume $x \not\equiv 2 \pmod 3$ since the graphs $CC_1(x, y)$ and $CC_1(y, x)$ are isomorphic.

First, when dealing with $\lambda \geq x + 5$, we can pick up, as V_λ , the λ first vertices of the ordering $\{u_1, \dots, u_x, a_1, b_1, c_1, d_1, e_1, v_y, \dots, v_1\}$ of $V(C)$ to get a partition of C into a traceable graph or $CC_1(x, y - (\lambda - (x + 6)))$ which is R-AP by the induction hypothesis, and a path. For $\lambda = x$, one can consider $V_\lambda = \{u_1, \dots, u_x\}$ so that the two induced graphs are traceable. Now, if $\lambda = x + 2$ or $\lambda = x + 3$, then we can choose $\{u_1, \dots, u_x, a_1, b_1\}$ or $\{u_1, \dots, u_x, a_1, c_1, d_1\}$, respectively, as V_λ , so that the two induced subgraphs are paths. Next, consider $\lambda = x + 4$. Then $V_\lambda = \{u_1, \dots, u_x, a_1, b_1, c_1, d_1\}$ yields a correct partition of C . Indeed, on the one hand, $C[V_\lambda]$ is a caterpillar $Cat(3, x + 1)$ which is R-AP since otherwise it would mean that $x \equiv 2 \pmod 3$, a contradiction. On the other hand, the graph $C[V_{n-\lambda}]$ is a path.

Now consider $\lambda = x + 1$. If $V_\lambda = \{u_1, \dots, u_x, a_1\}$ does not provide a satisfying partition of C , then $y \equiv 2 \pmod 3$ since $C[V_{n-\lambda}]$ is $Cat(3, y + 1)$ and is not R-AP. Consider now, as V_λ , the λ first vertices of the ordering $(v_1, \dots, v_y, e_1, b_1, d_1, c_1, a_1, u_x, \dots, u_1)$ of $V(C)$. If this choice of V_λ does not yield a correct partition of C once again, then it means that either $C[V_\lambda]$ is the caterpillar $Cat(3, y + 1)$, or a connected cycles graph $CC_1(x', y)$ with $x' \equiv 2 \pmod 3$. But then we get that either $x + 1 = y + 4$ or $x + 1 = y + 5 + x'$, respectively, which both imply that $x \equiv 2 \pmod 3$, a contradiction.

Finally consider every value $\lambda \in \{1, \dots, x - 1\}$. On the one hand, if $x - \lambda \not\equiv 2 \pmod 3$, then choose $V_\lambda = \{u_1, \dots, u_\lambda\}$ so that $C[V_\lambda]$ and $C[V_{n-\lambda}]$ are a path, and $CC_1(x - \lambda, y)$ which is R-AP by the induction hypothesis. On the other hand, i.e. $x - \lambda \equiv 2 \pmod 3$, we have $\lambda \not\equiv 0 \pmod 3$ since otherwise we would have $x \equiv 2 \pmod 3$. We can assume that $\lambda \notin \{y, y + 2, y + 3\}$, since otherwise we could deduce a correct partition of C as in the cases $\lambda \in \{x, x + 2, x + 3\}$, respectively. Then consider, as V_λ , the λ first vertices of $(v_1, \dots, v_y, e_1, b_1, d_1, c_1, a_1, u_x, \dots, u_1)$. If this choice of V_λ does not yield a correct partition of C , then $C[V_\lambda]$ is either a caterpillar $Cat(3, y + 1)$ which is not R-AP, or a graph $CC_1(x', y)$ with $x' \equiv 2 \pmod 3$. But note then that the first situation cannot occur because $\lambda \not\equiv 0 \pmod 3$. For the second situation, note that, because $\lambda \not\equiv 0 \pmod 3$, we have $y \not\equiv 2 \pmod 3$. Since we have $x', y < x$, the graph $CC_1(y, x')$ is actually R-AP by the induction hypothesis.

Case 2. Arbitrary k.

Let us now suppose that the result is true for every i up to $k - 1$, and let us prove it for k . Consider first $C = PCC_k(x)$ for consecutive values of $x \not\equiv 2 \pmod 3$. As we did before, to prove that C is R-AP we show that there exists a partition of $V(C)$ satisfying our conditions for every possible value of λ . One may choose V_λ as follows.

- If $\lambda \equiv 1 \pmod 3$, then one may consider, as V_λ , the first λ vertices of the ordering $(b_k, d_k, c_k, a_k, w_{k-1,k}, e_{k-1}, b_{k-1}, d_{k-1}, c_{k-1}, a_{k-1}, \dots, w_{1,2}, e_1, b_1, d_1, c_1, a_1, u_x, \dots, u_1)$ of $V(C)$. On the one hand, notice that $C[V_\lambda]$ is either a path, or covered by a R-AP caterpillar or a partial connected cycles graph $PCC_{k'}(x')$ with $k' \leq k - 1$ and $x' \not\equiv 2 \pmod 3$, which is R-AP by the induction hypothesis. On the other hand, observe that $C[V_{n-\lambda}]$ is either traceable, or spanned by a connected cycles graph $CC_{k''}(x, y)$ for some $k'' \leq k - 1$ and y , which is R-AP according to the induction hypothesis.

- If $\lambda \equiv 2 \pmod{3}$, then one can obtain similar partitions of C from the ordering $(d_k, c_k, b_k, a_k, w_{k-1,k}, e_{k-1}, d_{k-1}, c_{k-1}, b_{k-1}, a_{k-1}, \dots, w_{1,2}, e_1, d_1, c_1, b_1, a_1, u_x, \dots, u_1)$ of $V(C)$.
- Otherwise, if $\lambda \equiv 0 \pmod{3}$, then one has to consider as V_λ the first λ vertices of the ordering $(u_1, \dots, u_x, a_1, b_1, c_1, d_1, e_1, w_{1,2}, \dots, a_{k-1}, b_{k-1}, c_{k-1}, d_{k-1}, e_{k-1}, w_{k-1,k}, a_k, b_k, c_k, d_k)$ of $V(C)$ when $x \equiv 1 \pmod{3}$, or the ordering $(u_1, \dots, u_x, a_1, c_1, d_1, b_1, e_1, w_{1,2}, \dots, a_{k-1}, c_{k-1}, d_{k-1}, b_{k-1}, e_{k-1}, w_{k-1,k}, a_k, c_k, d_k, b_k)$ otherwise, i.e. when $x \equiv 0 \pmod{3}$. The two induced subgraphs $C[V_\lambda]$ and $C[V_{n-\lambda}]$ are then R-AP. Indeed, on the one hand, $C[V_\lambda]$ is either isomorphic to a path or spanned by a connected cycles graph $CC_{k'}(x, y)$ for $k' \leq k-1$ and some y . On the other hand, the subgraph $C[V_{n-\lambda}]$ is spanned by some $PCC_{k''}(x')$ graph with $k'' \leq k$ and $x' \not\equiv 2 \pmod{3}$.

To end up proving the claim, we have to show that $CC_k(x, y)$ is R-AP whenever $x \not\equiv 2 \pmod{3}$ or $y \not\equiv 2 \pmod{3}$. As for the base case, we show this by induction on $x + y$. Once again, we assume that $x \not\equiv 2 \pmod{3}$ for a given graph $C = CC_k(x, y)$.

For some $\lambda \in \{1, \dots, y\}$, one can consider $V_\lambda = \{v_1, \dots, v_\lambda\}$ so that C is partitioned into a path and $CC_k(x, y - \lambda)$ which is R-AP according to the induction hypothesis. When $\lambda = y + 1$, one can choose $V_\lambda = \{v_1, \dots, v_y, e_k\}$ so that C is partitioned into a path and a partial connected cycles graph which is R-AP by the induction hypothesis since $x \not\equiv 2 \pmod{3}$. For other values of λ , one may choose V_λ as follows.

- If $\lambda \equiv 0 \pmod{3}$, one can consider, as V_λ , the λ first vertices from the ordering $(u_1, \dots, u_x, a_1, b_1, c_1, d_1, e_1, w_{1,2}, \dots, w_{k-1,k}, a_k, b_k, c_k, d_k, e_k, v_y, \dots, v_1)$ of $V(C)$ when $x \equiv 1 \pmod{3}$, from $(u_1, \dots, u_x, a_1, c_1, d_1, b_1, e_1, w_{1,2}, \dots, w_{k-1,k}, a_k, c_k, d_k, b_k, e_k, v_y, \dots, v_1)$ otherwise, i.e. when $x \equiv 0 \pmod{3}$. The two induced subgraphs are then R-AP since they are traceable or isomorphic to connected cycles graphs which are R-AP according to the induction hypotheses.
- If $\lambda \not\equiv 0 \pmod{3}$ and $\lambda - (y + 1) \equiv 0 \pmod{3}$, then one can consider the λ first vertices of the ordering $(v_1, \dots, v_y, e_k, b_k, d_k, c_k, a_k, w_{k-1,k}, \dots, e_1, b_1, d_1, c_1, a_1, u_x, \dots, u_1)$ of $V(C)$. For each such partition, we get, on the one hand, that $C[V_\lambda]$ is either a path, a R-AP caterpillar, or a R-AP (partial) connected cycles graph. In particular, note that when $C[V_\lambda]$ is a caterpillar or partial connected cycles graph, then this graph is R-AP since $y \not\equiv 2 \pmod{3}$ because of the assumptions on λ . On the other hand, the graph $C[V_{n-\lambda}]$ is either a path, or a (partial) connected cycles graph which is R-AP by the induction hypothesis.
- If $\lambda \not\equiv 0 \pmod{3}$ and $\lambda - (y + 1) \equiv 1 \pmod{3}$, then one may pick up, as V_λ , the λ first vertices from the ordering given to deal with the previous case. This choice of V_λ makes, on the one hand, $C[V_\lambda]$ being spanned by either a path, or $CC_{k'}(x', y)$ where $k' \leq k-1$ and $x' \not\equiv 2 \pmod{3}$ which is R-AP by the induction hypothesis. On the other hand, $C[V_{n-\lambda}]$ is a path, or is spanned by some graph $CC_{k''}(x, y')$ for $k'' \leq k-1$ and some y' which is R-AP, again by the induction hypothesis.
- Otherwise, if $\lambda \not\equiv 0 \pmod{3}$ and $\lambda - (y + 1) \equiv 2 \pmod{3}$, then some similar partitions of C may be obtained from the ordering $(v_1, \dots, v_y, e_k, d_k, c_k, b_k, a_k, w_{k-1,k}, \dots, w_{1,2}, e_1, d_1, c_1, b_1, a_1, u_x, \dots, u_1)$ of $V(C)$. \square

Note that Lemma 3.4 provides a full characterization of R-AP (partial) connected cycles graphs since every such graph whose parameters do not satisfy this lemma is not R-AP. To be convinced of that fact, one just has to consider successive partitions of such a graph for $\lambda = 3$.

We finally deduce Theorem 3.1 as a corollary of Lemma 3.4.

Proof of Theorem 3.1. We have $LP(CC_{c+1}(1, 1)) = |V(CC_{c+1}(1, 1))| - (c + 1)$ for every $c \geq 1$. Therefore, for every possible value of c , we have a graph showing that c does not contradict the claim. \square

Finally notice that by adding the edge u_1v_1 to any connected cycles graph $CC_k(1, 1)$, we get a 2-connected graph which is R-AP according to Remark 2.1 and whose longest path has order $LP(CC_k(1, 1)) + 1$. Therefore, Theorem 3.1 is also true when restricted to 2-connected graphs.

4. LONGEST PATH AND MULTIPLICATIVE FACTOR

The graph $CC_k(1, 1)$ has order $n = 6k + 1$ while its longest path has order $n - k$ for every $k \geq 1$. Thus, even if the connected cycles graphs confirm that the order of the longest path in a R-AP n -graph is not constantly lower than n up to an additive factor, they do not reject the strong relationship between the properties of being R-AP and traceable. We now discuss how to create this relationship thanks to a multiplicative factor.

Question 4.1. *What is the biggest $c < 1$ such that $LP(G) \geq n \cdot c$ for every R-AP n -graph G ?*

Regarding the connected cycles graphs, we get that $c \leq \frac{5}{6}$. In this section, we deduce a better upper bound on c thanks to the following graph construction.

Definition 4.2. Let $k, k' \geq 1$ be two positive integers. The *urchin* $W(k, k')$ is the graph obtained as follows.

- Let A, B, C be three sets of k, k and k' distinct vertices, respectively.
- Add a perfect matching between the vertices of A and B .
- Add all possible edges between distinct vertices in $B \cup C$.

This construction is illustrated in Figure 3. Note that the urchin $W(k, k)$ has order $3k$ while its longest path has order $2k + 2$. We then get that $LP(W(k, k))/n$ tends to $\frac{2}{3}$ as k grows to infinity. In what follows, we show that any urchin $W(k, k)$ is R-AP, and thus that the following holds regarding Question 4.1.

Theorem 4.3. *Regarding Question 4.1, we have $c \leq \frac{2}{3}$.*

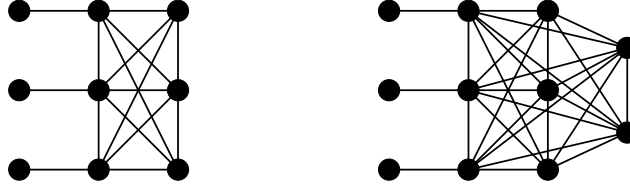


Fig. 3. The urchins $W(3,3)$ and $W(3,5)$

We prove that an urchin $W(k, k')$ is R-AP for some values of k and k' .

Lemma 4.4. *The urchin $W(k, k')$ is R-AP for every $k \geq 2$ and $k' \geq k - 2$.*

Proof. We introduce some terminology to deal with the vertices of any urchin $W(k, k')$. The vertices of A are denoted u_1, \dots, u_k , and those of B are denoted v_1, \dots, v_k in such a way that $u_i v_i$ is an edge for every $i \in \{1, \dots, k\}$. The vertices of C are denoted $w_1, \dots, w_{k'}$ arbitrarily.

The claim is proved by induction on both k and k' . As a base case, note that every urchin $W(2, k')$ is traceable, and thus R-AP by Remark 2.2. Suppose now that $W(i, i')$ is R-AP for every i up to $k - 1$ and $i' \geq i - 2$. We now prove that the urchin n -graph $W = W(k, k')$ is R-AP for every $k' \geq k - 2$. For this purpose, we show, for every value of $\lambda \in \{1, \dots, \lfloor \frac{n}{2} \rfloor\}$, that $V(W)$ can be partitioned into two parts V_λ and $V_{n-\lambda}$ inducing R-AP graphs on λ and $n - \lambda$ vertices, respectively.

We first deal with the easy cases, i.e. $\lambda \in \{1, 2, 3\}$. For $\lambda = 1$, consider $V_\lambda = \{u_1\}$ so that the two induced subgraphs are K_1 and $W(k - 1, k' + 1)$. Since $k' \geq k - 2$, this subgraph is R-AP by the induction hypothesis. For $\lambda = 2$, let $V_\lambda = \{u_1, v_1\}$. The two induced subgraphs then are K_2 and $W(k - 1, k')$, which is R-AP for the same reason as the previous case. Now, for $\lambda = 3$, choose $V_\lambda = \{u_1, v_1, w_1\}$. The two induced subgraphs then are a path, and the urchin $W(k - 1, k' - 1)$ which is R-AP, again by the induction hypothesis.

We now deal with the remaining values of λ , i.e. $\lambda \geq 4$. The part V_λ is obtained by choosing two disjoint sets V'_λ and V''_λ , and then considering their union. On the one hand, in the case where $\lambda \equiv 1 \pmod 3$, let $x = \lfloor \frac{\lambda-4}{3} \rfloor$. Clearly, x is an integer. First, let $V'_\lambda = \emptyset$ if $x = 0$, or $V'_\lambda = \bigcup_{i=1}^x \{u_i, v_i, w_i\}$ otherwise. Then set $V''_\lambda = \{v_{x+1}, u_{x+1}, v_{x+2}, u_{x+2}\}$. The two induced subgraphs then are a path or $W(x + 2, x)$, and $W(k - (x + 2), k' - (x - 2))$, which are R-AP by the induction hypothesis since $k' \geq k - 2$.

On the other hand, i.e. $\lambda \not\equiv 1 \pmod 3$, let $x = \lfloor \frac{\lambda}{3} \rfloor$ and $y \equiv \lambda \pmod 3$ with $y \in \{0, 2\}$. Then, let $V'_\lambda = \bigcup_{i=1}^x \{u_i, v_i, w_i\}$. The strategy for choosing V''_λ depends on whether $y = 0$ or $y = 2$.

- $y = 0$. Choose $V''_\lambda = \emptyset$. In this situation, the two induced subgraphs are $W(x, x)$ and $W(k - x, k' - x)$ which are R-AP by the induction hypothesis since $k' \geq k - 2$.
- $y = 2$. Let $V''_\lambda = \{v_{x+1}, u_{x+1}\}$. The two induced subgraphs then are $W(x + 1, x)$ and $W(k - (x + 1), k' - x)$, which are R-AP according to the induction hypothesis. \square

Theorem 4.3 follows as a corollary of Lemma 4.4. Note that Lemma 4.4 is tight in the sense that urchins $W(k, k - x)$ with $x \geq 3$ are not R-AP since such a graph W

cannot be partitioned as requested for $\lambda = 3$. Indeed, as a set V_λ with size 3 inducing a R-AP subgraph of W , one has to consider, following the terminology introduced in the proof of Lemma 4.4, a part of the form $\{u_i, v_i, w_j\}$ or $\{w_i, w_j, w_\ell\}$. After having successively picked several sets with size 3 off W , one necessarily gets an urchin $W(k', 0)$ with $k' \geq 3$. Such a graph is clearly not partitionable for $\lambda = 3$ once again.

We can strengthen Theorem 4.3 as follows. Let $W = W(k, k')$ be a R-AP urchin. Observe that by adding the edges u_1u_2, \dots, u_1u_k to W , we get a 2-connected graph W_2 which is R-AP by Remark 2.1. By then adding the edges u_2u_3, \dots, u_2u_k to W_2 , we get another R-AP graph W_3 which is 3-connected. By repeating this procedure as many times as needed, we get an ℓ -connected R-AP graph W_ℓ for any value of ℓ assuming k and k' are big enough. Note that we have $LP(W_i) = LP(W) + 2i$, and thus that $LP(W_i)/LP(W)$ tends to 1 as k grows to infinity. Therefore, the statement of Theorem 4.3 is also true when restricted to ℓ -connected R-AP graphs, no matter what is the value ℓ .

Theorem 4.5. *Theorem 4.3 is also true when Question 4.1 is restricted to R-AP graphs of arbitrary connectivity.*

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